# Cryptography and Security Exam 

## 2nd Exam

20.2.2008

## Crypto Part

## 1 AES-Hashing

In this exercise we consider a special hash function $H$ defined as follows. To hash a message $m$ with a length multiple of 256 bits, we split it into blocks of 256 bits $m_{1}, \ldots, m_{b}$. Then, we compute the encryption of $i$ with key $m_{i}$ using AES for $i=1, \ldots, b$ and XOR them all together. We define

$$
H\left(m_{1}\|\cdots\| m_{b}\right)=\bigoplus_{i=1}^{b} \mathrm{AES}_{m_{i}}(i)
$$

1. What is the length of the digest?

Ideally, what should be the complexity of the best collision attack on $H$ ?
Ideally, what should be the complexity of the best preimage attach on $H$ ?
2. Derive a collision attack to find two messages $m$ and $m^{\prime}$ of length 256 bits with same digest. What is its complexity?
3. Derive a preimage attack to find a preimage of the digest 0 and finding a message of length 512 bits.
What is its complexity?
4. Derive a second preimage attack finding a message of length 512 bits for any first preimage.

What is its complexity?
5. Let $m$ and $m^{\prime}$ be two messages of same bitlength $256 b$ for an integer $b$. Let $m=m_{1}\|\cdots\| m_{b}$ and $m^{\prime}=m_{1}^{\prime}\|\cdots\| m_{b}^{\prime}$ be the decomposition into 256-bit blocks. We assume that $m$ and $m^{\prime}$ are selected such that $m_{i} \neq m_{i}^{\prime}$ for $i=1, \ldots, b$. Let $u_{i}=\mathrm{AES}_{m_{i}}(i) \oplus \mathrm{AES}_{m_{i}^{\prime}}(i)$.
How large should $b$ be so that with high probability, for any $y$ there exists a subset $I$ of $\{1, \ldots, b\}$ such that $y=\bigoplus_{i \in I} u_{i}$ ?
By selecting $b$ this way, derive a preimage attack which finds a message of length $256 b$ bits for any digest $h$. (Hint: set $y=h \oplus H(m)$.)
What is its complexity?

## 2 Modulo 11 Diffie-Hellman

1. Let $d_{n-1} \ldots d_{1} d_{0}$ be the decimal expansion of an integer $N$, i.e. $d_{i} \in\{0,1, \ldots, 9\}$ and

$$
N=\sum_{i=0}^{n-1} 10^{i} \times d_{i}
$$

Show that $N \equiv d_{0}-d_{1}+\cdots+(-1)^{n-1} d_{n-1} \quad(\bmod 11)$.
Deduce an algorithm to reduce an integer modulo 11 by mental computing.
2. What is the order of the $\mathbf{Z}_{11}^{*}$ group?

Show that 2 is a generator of $\mathbf{Z}_{11}^{*}$.
What is the order of 3 in $\mathbf{Z}_{11}^{*}$ ?
3. Consider the Diffie-Hellman protocol with prime number $p=11$ and generator $g=2$. Alice picks an exponent $x=9$, sends $X=g^{x} \bmod p$ to Bob and gets $Y=8$ from him.
Compute $X$.
Compute the Diffie-Hellman key $K$.

## 3 Modulo 1111 RSA

1. Let $d_{n-1} \ldots d_{1} d_{0}$ be the basis-100 expansion of an integer $N$, i.e. $d_{i} \in\{0,1, \ldots, 99\}$ and

$$
N=\sum_{i=0}^{n-1} 100^{i} \times d_{i}
$$

Show that $N \equiv d_{0}-d_{1}+\cdots+(-1)^{n-1} d_{n-1} \quad(\bmod 101)$.
Deduce an algorithm to reduce an integer modulo 101 by mental computing.
2. With the same notations, show that $N \equiv \sum_{i} d_{i}(\bmod 11)$.

Deduce an algorithm to reduce an integer modulo 11 by mental computing.
3. Let $a$ and $b$ be arbitrary integers and let $N=(6 \times 101 \times a+46 \times 11 \times b) \bmod 1111$.

Show that $N \equiv a \quad(\bmod 11)$ and $N \equiv b \quad(\bmod 101)$.
Show that $N$ is the unique integer with this property in the $[0,1110]$ interval.
4. Consider RSA signatures with public key $N=1111$ and $e=3$.

Compute the secret key $d$.
Compute the signature $y$ of the message $x=2$.

