

Family Name:
First Name:
Section:

# Security and Cryptography

Midterm Exam (Solutions)

October  $30^{\text{th}}$ , 2008

Duration: 1 hour 45 min

# Instructions

Documents are not allowed apart from linguistic dictionaries.
Electronic devices (including calculators ) are not allowed.
Answers must be written on the exercises sheet.
This exam contains 3 independent exercises.
Answers can be either in French or English.
Questions of any kind will certainly not be answered.
Potential errors in these sheets are part of the exam.
You have to put your full name on the first page and have all pages stapled.

#### 1 Square roots of 53 modulo 221

The purpose of this exercise is to solve in  $\mathbf{Z}_n$  the equation

$$x^2 \equiv a \pmod{n}$$

with n = 221 and a = 53.

1. Let n = pq be the factorization of n into prime numbers where p is the smallest one. Compute p and q.

Since  $13 \times 17 = 221$  we have p = 13 and q = 17.

2. Solve in  $\mathbf{Z}_p$  the equation  $x^2 \equiv a$ .

We have a mod  $p = 53 \mod 13 = 1$ . Since p is a prime number the only square roots of +1 are +1 and -1. That is, x mod p is either 1 or 12.

3. Solve in  $\mathbf{Z}_q$  the equation  $x^2 \equiv a$ .

We have a mod q = 2. We have  $a^{\frac{q-1}{2}} = 2^8 \pmod{17}$ . This is  $(-1)^4 \mod 16$  which is 1, so a is a quadratic residue. To find a square root, we can try by exhaustive search as q is pretty small. The list of squares modulo 17 is  $1, 4, 9, 16, 8, 2, \ldots$  so 6 is a quare root of 1. The second square root is  $-6 = 11 \pmod{17}$ .

4. Reduce  $\alpha = 170 \mod p$  and modulo q.

We compute 170 mod p = 1 and 170 mod q = 0

5. Reduce  $\beta = 1 - \alpha$  modulo p and modulo q.

Clearly,  $(1 - 170) \mod p = 0$  and  $(1 - 170) \mod q = 1$ 

6. Given arbitrary u and v, reduce  $u\alpha + v\beta$  modulo p and q.

Clearly,  $(u\alpha + v\beta) \mod p = u \mod p$  and  $(u\alpha + v\beta) \mod q = v \mod q$ 

7. List all roots in  $\mathbf{Z}_n$  of the equation  $x^2 \equiv a$ .

We use the Chinese remainder theorem. We have  $x^2 \mod n = a \iff x^2 \mod p = a \mod p$  and  $x^2 \mod q = a \mod q$ . The right hand side suggests 4 solutions such that  $x \mod p \in \{1, 12\}$  and  $x \mod q \in \{6, 11\}$ . Using the previous question we compute x by

 $\begin{array}{rcl} 1 \times \alpha + 6 \times \beta &=& 170 \times (1-6) + 6 \\ 1 \times \alpha - 6 \times \beta &=& 170 \times (1+6) - 6 \\ -1 \times \alpha + 6 \times \beta &=& 170 \times (-1-6) + 6 \\ -1 \times \alpha - 6 \times \beta &=& 170 \times (-1+6) - 6 \end{array}$ 

We do the computation modulo 221 and obtain  $x \in \{40, 79, 142, 181\}$ .

### 2 RSA with exponent 3

In this exercise we consider an RSA modulus n = pq where p and q are large prime numbers (here, by "large" we mean at least equal to 5). We consider a valid RSA exponent e for RSA.

1. Show that neither  $p \mod 3$  nor  $q \mod 3$  can be equal to 0.

Since p and q are large, they are larger than 3. Since they are prime, they are not divisible by 3.

2. Under which condition e is a valid exponent for a modulus n?

e is a valid exponent iff it is coprime with  $\varphi(n) = (p-1)(q-1)$ .

From now on, we will assume that e = 3.

3. Show that neither p-1 nor q-1 can be multiples of 3.

If e = 3 is a valid exponent, then (p-1)(q-1) is coprime with 3, which means that it is not divisible by 3. Therefore, neither p-1 nor q-1 can be divisible by 3.

4. Deduce that  $p \mod 3 = q \mod 3 = 2$ .

Due to the previous questions,  $p \mod 3$  is neither 0 nor 1 so it must be 2. The same holds for q.

5. What is the value of  $n \mod 3$ ?

We have  $n = p \times q = 2^2 = 1 \pmod{3}$ .

6. For any digits  $d_0, \ldots, d_{\ell-1}$ , show that

$$\left(\sum_{i=0}^{\ell-1} d_i 10^i\right) \mod 3 = \left(\sum_{i=0}^{\ell-1} (d_i \bmod 3)\right) \mod 3$$

This directly comes from  $10 \mod 3 = 1$ .

7. Show that e = 3 is not a valid RSA exponent for the following RSA modulus:

 $n = 777\,575\,993$ 

From the previous question we have  $n \mod 3 = (1+1+1+2+1+2+0+0+0) \mod 3 = 2$  which is not equal to 1. So, either n is not a product of two primes or 3 is not a valid exponent. In any case, (n, 3) is not a valid RSA public key.

## **3** Computation in GF(16)

Let us consider the polynomial  $P(x) = x^4 + x + 1$  in  $\mathbb{Z}_2[x]$ .

1. Show that P has no root in  $\mathbb{Z}_2$ .

We have P(0) = 1 and P(1) = 1 so it has no root in  $\mathbb{Z}_2$ .

2. Deduce that P has no factor of degree 1 in  $\mathbb{Z}_2[x]$ .

Having a factor of degree 1 is equivalent to having a root. So, P has no factor of degree 1.

3. Enumerate all polynomials of degree 2 in  $\mathbf{Z}_2[x]$  and identify the one Q(x) which is irreducible.

We have  $x^2$ ,  $x^2 + 1$ ,  $x^2 + x$ ,  $x^2 + x + 1$ . We can check that all have roots, except  $x^2 + x + 1$ . So, only  $x^2 + x + 1$  remains as a candidate for being irreducible. Since it has degree 2, having no factor of degree 1 is enough to guaranty irreducibility. Hence,  $Q(x) = x^2 + x + 1$  is the only irreducible polynomial of degree 2.

4. Show that Q(x) does not divide P(x).

We have  $P(x) = x^4 + x + 1 = (x^2 + x) \times Q(x) + 1$  so P(x) is not divisible by Q(x).

5. Deduce that P(X) is irreducible.

P(x) has degree 4 and no factor of degree 1. Thus, either it has two irreducible factors of degree 2 or it is irreducible. Since Q(x) is the only irreducible polynomial of degree 2 and is not a factor of P(x), P(x) must be irreducible.

6. We define

$$\mathsf{GF}(16) \leftrightarrow \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F\}$$

where an hexadecimal  $u = \alpha 2^0 + \beta \times 2^1 + \gamma \times 2^2 + \delta \times 2^3$  with  $\alpha, \beta, \gamma, \delta \in \{0, 1\}$  is considered to represent the polynomial

$$\alpha + \beta x + \gamma x^2 + \delta x^3$$
 in GF(16)

Those polynomials in  $\mathbf{Z}_2[x]$  are taken modulo P(x).

(a) What is the GF(16)-sum of 6 and A?

We have  $6 = 2^2 + 2^1$  so it represents  $x^2 + x$ . We have  $A = 2^3 + 2^1$  so it represents  $x^3 + x$ . Thus, the sum is  $x^3 + x^2$  which is represented by  $2^3 + 2^2 = C$ :  $6 \boxplus A = C$ .

(b) What is the GF(16)-multiplication of 6 and 1?

Since 1 represents 1, the multiplication by 1 is trivial:  $6 \boxtimes 1 = 6$ .

(c) What is the GF(16)-multiplication of 6 and 2?

Since 2 represents x, the multiplication by 2 can be done by shifting bits (if no carry). 6 represents  $x^2 + x$  which is shifted to  $x^3 + x^2$ , represented by C:  $6 \boxtimes 2 = C$ .

(d) What is the GF(16)-multiplication of 6 and 3?

We can check that  $3 = 2 \boxplus 1$  (indeed). We have  $6 \boxtimes 3 = (6 \boxtimes 2) \boxplus (6 \boxtimes 1) = C \boxplus 6 = A$ .

(e) What is the GF(16)-inverse of 2?

We have  $x^4 + x + 1 = x \times (x^3 + 1) + 1$  so  $0 = (2 \boxtimes 9) \boxplus 1$  which can also writes  $2 \boxtimes 9 = 1$ : 9 is the inverse of 2.